## Monthly Problem

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February 14, 2007

11007. Proposed by Western Maryland College Problems Group, Westminster, MY. Let  $\langle \rangle$  denote Eulerian numbers, and let  $\{ \}$  denote Stirling numbers of the second kind. Show that

$$\sum_{j=1}^{n} 2^{j-1} \left\langle {n \atop j} \right\rangle = \sum_{j=1}^{n} j! \left\{ {n \atop j} \right\}.$$
<sup>(1)</sup>

Solution: Let  $\mathbb{N}_k$  denote the set of the first k positive integers. Then each term in the summation on the right-hand-side of (1) counts the number of onto functions  $f : \mathbb{N}_n \twoheadrightarrow \mathbb{N}_j$ . This follows because the Stirling number of the second kind,  ${n \atop j}$ , counts the number of unordered partitions of a set of cardinality n into j non-empty classes. So  $j! {n \atop j}$  gives the number of such ordered partitions, which is the same as the number of onto functions  $f : \mathbb{N}_n \twoheadrightarrow \mathbb{N}_j$ .

In order to establish the identity, we need a connection between the number of onto functions and the Eulerian numbers. This is given in the following identity which appears in [1] using a different notation where its proof is left as an exercise.

Lemma 1.

$$j! \begin{Bmatrix} n \\ j \end{Bmatrix} = \sum_{k=0}^{n-1} \begin{Bmatrix} n \\ k+1 \end{Bmatrix} \begin{Bmatrix} k \\ n-j \end{Bmatrix}.$$
(2)

Proof (of lemma). Since the left-hand side of (2) counts the number of onto functions from  $\mathbb{N}_n \twoheadrightarrow \mathbb{N}_j$ , we show the right-hand side counts these functions as well.

So let  $f : \mathbb{N}_n \twoheadrightarrow \mathbb{N}_j$  be onto. Then f can be represented as an n-tuple of integers  $(a_1, a_2, \dots, a_n)$  where  $1 \le a_i \le j$  for each i and where  $\{a_i | 1 \le i \le j\} = \mathbb{N}_j$  (since f is onto). Now rearrange the permutation so all the elements of  $f^{-1}(1)$  occur first and in increasing order, then the elements of  $f^{-1}(2)$  occur next arranged in increasing order, etc., to obtain  $(a_{i_1}, a_{i_2}, \dots, a_{i_n})$ . Then  $(i_1, i_2, \dots, i_n)$  is a permutation on  $\mathbb{N}_n$  with j - 1 or fewer descents (where a descent occurs when one number in a permutation is less that its predecessor).

Now, given any n-permutation with j - 1 or fewer descents, we use the descents as barriers to induce a partition on the n numbers. We add j - 1 - k additional barriers, where k is the number of descents, to construct an onto function  $f : \mathbb{N}_n \twoheadrightarrow \mathbb{N}_j$ . This can be done by choosing where the j - 1 - k barriers go from

the n-1-k available positions. This can be done in  $\binom{n-1-k}{j-1-k} = \binom{n-1-k}{n-j}$  ways. Summing over all possible k gives the number of onto functions.

$$\sum_{k=0}^{n-1} {\binom{n}{k+1} \binom{n-1-k}{n-j}} = \sum_{k=0}^{n-1} {\binom{n}{k+1} \binom{k}{n-j}}$$

An example will serve to illustrate this counting method. A similar example is found in [2]. If we are counting the number of onto functions  $f : \mathbb{N}_9 \to \mathbb{N}_6$  and consider the permutation on  $\mathbb{N}_9$  given by 135274698, how many onto functions does this permutation correspond to with this counting method? If we write the permutation showing the descents, we get  $135 \downarrow 27 \downarrow 469 \downarrow 8$ . To define an onto function to  $\mathbb{N}_6$ , we must insert 2 additional barriers between adjacent numbers to create a total of 5 barriers to produce a partition into 6 ordered classes. Considering the available locations for additional barriers (denoted by  $\sqcup$ ),  $1 \sqcup 3 \sqcup 5 \downarrow 2 \sqcup 7 \downarrow 4 \sqcup 6 \sqcup 9 \downarrow 8$ , we must choose 2 of the 5 without regard to order in any of  $\binom{5}{2}$  ways.  $\Box$ 

Using the lemma, we have

$$\begin{split} \sum_{j=1}^{n} j! \left\{ \begin{array}{l} n \\ j \end{array} \right\} &= \sum_{j=1}^{n} \sum_{k=0}^{n-1} \left\langle \begin{array}{l} n \\ k+1 \end{array} \right\rangle \begin{pmatrix} k \\ n-j \end{pmatrix} \\ &= \sum_{j=1}^{n} \sum_{k=1}^{n} \left\langle \begin{array}{l} n \\ k \end{array} \right\rangle \begin{pmatrix} k-1 \\ n-j \end{pmatrix} \\ &= \sum_{k=1}^{n} \sum_{j=1}^{n} \left\langle \begin{array}{l} n \\ k \end{array} \right\rangle \begin{pmatrix} k-1 \\ n-j \end{pmatrix} \\ &= \sum_{k=1}^{n} \left\langle \begin{array}{l} n \\ k \end{array} \right\rangle \sum_{j=1}^{n} \begin{pmatrix} k-1 \\ n-j \end{pmatrix} \\ &= \sum_{k=1}^{n} \left\langle \begin{array}{l} n \\ k \end{array} \right\rangle 2^{k-1}. \end{split}$$

## References

- R.L. Graham, D.E. Knuth, and O. Patashnik, Concrete Mathematics, Second Edition, Addison-Wesley, Reading, MA, 1994.
- [2] Donald E. Knuth, The Art of Computer Programming, Volume 3: Sorting and Searching, Second Edition, 1998, Addison-Wesley.